## The Golden Digits International Contest

6th Edition, November 2024



**Problem 1.** Find all functions  $f : \mathbb{R} \to \mathbb{R}$ , such that for any real numbers x, y with  $y \neq 0$  we have:

$$f(f(x) + y)f\left(\frac{1}{y}\right) = xf\left(\frac{1}{y}\right) + 1$$

MARIUS CERLAT

**Solution:** Quite clearly,  $f(\mathbb{R}^*) \neq \{0\}$ , so there is some  $y \in \mathbb{R}^*$  with  $f\left(\frac{1}{y}\right) \neq 0$ . Substituting this y in the hypothesis yields that f is injective.

Again, for some  $y \in \mathbb{R}^*$  such that  $f\left(\frac{1}{y}\right) \neq 0$ , take  $x = -\frac{1}{f\left(\frac{1}{y}\right)}$ . This gives

$$f\left(y+f\left(-\frac{1}{f\left(\frac{1}{y}\right)}\right)\right) = 0$$

Therefore, there is some t with f(t) = 0. Putting x = t in the hypothesis will return

$$f(y)f\left(\frac{1}{y}\right) = tf\left(\frac{1}{y}\right) + 1$$

and since f is injective, swapping y with 1/y will give t = 0 and  $f(y)f\left(\frac{1}{y}\right) = 1$ .

**Claim:** f(x + y) = f(x) + f(y)

**Proof:** Multiplying the hypothesis by f(y) gives:

$$f(f(x) + y) = x + f(y) \tag{1}$$

Setting  $y \to f(y)$  in (1) yields f(f(x) + f(y)) = x + f(f(y)) = f(f(x)) + y, for all nonzero x, and y. Thus, there is a constant c such that f(f(x)) = x + c for all  $x \neq 0$ . Setting x = -c in this relation gives  $f(f(-c)) = 0 \Rightarrow c = 0 \Rightarrow f(f(x)) = x$ . Now, setting  $x \to f(x)$  in (1) gives the desired conclusion.

Notice that for any  $x \in \mathbb{R}$  with  $|x| \ge 2$ , there is some  $y \in \mathbb{R}^*$  s.t.  $x = y + \frac{1}{y}$ . Thus, by AM-GM,

$$|f(x)| = \left| f\left(y + \frac{1}{y}\right) \right| = \left| f(y) + f\left(\frac{1}{y}\right) \right| \ge 2.$$

Then, if  $|y| \leq \frac{1}{2}$   $(y \neq 0)$ ,  $|f(y)| = \frac{1}{|f(\frac{1}{y})|} \leq \frac{1}{2}$ , so f is bounded on  $\left(-\frac{1}{2}, \frac{1}{2}\right)$ , implying that f is linear. A quick check gives f(x) = x and f(x) = -x as the only solutions.

**Problem 2.** Vadim and Marian play a game. Starting with Vadim, they take turns eliminating exactly one edge from a complete graph with 2024 vertices. The first player to make a move that leaves no cycles loses. Determine who has a winning strategy.

## PAVEL CIUREA

## Solution: We show that Marian wins.

After Marian moves, the number of edges left uncut is even, so for him to lose the last cycle must be an odd one.

t := 2024/4. It is well known that if the number of uncut edges is  $\geq 4t$ , there is at least one cycle left. So, if Marian can make the graph bipartite before there are < 4t edges left he wins. We show that he can do so.

Marian splits the vertices in two sets,  $V_1$  and  $V_2$ , with  $|V_1| = |V_2| = 2t$ . There are  $2\binom{2t}{2}$  edges uv with  $u, v \in V_i$  with  $i \in \{1, 2\}$  (call such an edge good), and before there are < 4t edges left uncut, Marian cuts  $\binom{4t}{2} - 4t/2 = 4t^2 - 3t$  edges. Thus, it is enough for him to choose the sets such that Vadim cuts at least  $2\binom{2t}{2} - (4t^2 - 3t) = t$  good edges and he only cuts good ones.

At a certain point, we say a vertex u is *used* if there is an edge uv that has been cut.

**Lemma:** Marian can play so that after 2k - 1 turns there are at most 2k + 1 used vertices.

**Proof:** We show the statement by induction. The statement trivially holds for k = 2. We suppose the statement holds for k and show that it holds for k + 1.

Between the used vertices there are  $\binom{2k+1}{2} > 2k-1$  edges, so Marian can cut an edge between two already used vertices (if there are u < 2k+1 used vertices, then we consider 2k+1-u other vertices to be used), and since at a turn a player can make at most two new vertices used the statement is also true for k+1 so the lemma is proven.

Therefore, after 2t - 3 turns there are at most 2t - 1 used vertices. Marian adds all these vertices to  $V_1$ , then adds random vertices until  $|V_1|$  is 2t - 1 and cuts an edge between two vertices in  $V_1$ . On the next turn Vadim cuts an edge uv. If  $u \in V_1$  then Marian adds v in  $V_1$  (if  $v \in V_1$  he adds a random vertex), and if  $u, v \notin V_1$  he adds u and v to  $V_2$  and adds a random vertex to  $V_1$ .

This way, after 2t - 1 turns Vadim cut t good edges and Marian cut only good edges, so the constructed sets satisfy the desired conditions and thus Marian wins.

**Problem 3.** Prove that there exist infinitely many d such that we can find a polynomial P of degree d with integer coefficients and  $N \in \mathbb{N}$  such that for all integers x > N and any prime p we have:

$$v_p(P(x)^3 + 3P(x)^2 - 3) < \frac{d \cdot \log(x)}{2024^{2024}}$$

where  $\log(x)$  denotes the natural logarithm and  $v_p(n)$  denotes the largest number k such that  $p^k \mid n$ .

## MARIUS CERLAT

**Solution:** Let  $Q(x) = x^3 + 3x^2 - 3$ , and define  $Q(x - 1) = x^3 - 3x - 1$ . Inspired by the identity

$$\left(a+\frac{1}{a}\right)^3 - 3\left(a+\frac{1}{a}\right) = a^3 + \frac{1}{a^3},$$

we find that

$$Q\left(x + \frac{1}{x} - 1\right) = x^3 + \frac{1}{x^3} - 1.$$

Now, define  $P_n(x)$  as the polynomial with integer coefficients that satisfies

$$P_n\left(x+\frac{1}{x}-1\right) = x^n + \frac{1}{x^n} - 1.$$

Its existence and uniqueness can be easily proven by induction. We can observe that

$$P_n(P_m(x+\frac{1}{x}-1)) = x^{mn} + \frac{1}{x^{mn}} - 1 = P_m(P_n(x+\frac{1}{x}-1)).$$

Now, by substituting  $P_n(x)$  instead of P(x) in the initial condition, we find that our polynomial satisfies

$$P_3(P_n(x)) = P_n(P_3(x)).$$

**Claim:** For any positive constant C, there are infinitely many natural numbers n such that the degree of any irreducible polynomial dividing  $P_n(x)$  is at most Cn.

**Subclaim 1:** For any positive constant C, there are infinitely many natural numbers n such that the degree of any irreducible polynomial dividing  $x^{2n} - x^n + 1$  is at most Cn.

**Proof:** The polynomial  $x^{2n} - x^n + 1$  divides  $x^{3n} + 1$ , which divides

$$x^{6n} - 1 = \prod_{d|6n} \Phi_d(x),$$

where  $\Phi_d(x)$  is the *d*-th cyclotomic polynomial. It is well-known that  $\deg(\Phi_d) = \varphi(d)$ , where  $\varphi$  is Euler's totient function. Hence, for any divisor *d* of 6n,

$$\frac{\deg(\Phi_d)}{n} \le \frac{\varphi(6n)}{n} = 6 \prod_{p_i|6n} \left(\frac{p_i - 1}{p_i}\right),$$

where  $p_i$  are the prime factors of 6n. It is known that

$$\prod\left(\frac{p-1}{p}\right)$$

converges to 0. Therefore, by taking n to be divisible by primes  $p_1, p_2, \ldots, p_k$  such that

$$\prod \frac{p_i - 1}{p_i} < \frac{C}{6},$$

the claim is proved.

**Subclaim 2:** If a polynomial Q(x) is irreducible, then  $Q_R(x) = x^{\deg(Q)}Q\left(x + \frac{1}{x}\right)$  has at most two irreducible factors.

**Proof:** Assume  $Q_R(x) = M(x)N(x)$ , where M(x) is an irreducible polynomial. Our goal is to prove that N(x) must also be irreducible. We have

$$M(x)N(x) = x^{\deg(Q)}Q\left(x + \frac{1}{x}\right) = x^{2\deg(Q)}M\left(\frac{1}{x}\right)N\left(\frac{1}{x}\right).$$

This implies that the polynomial

$$\bar{M}(x) = x^{\deg(M)} M\left(\frac{1}{x}\right)$$

is irreducible. Now, we consider two cases:

1. If M(x) divides M(x), then since both polynomials have the same degree, we conclude that M(x) is simply a scaled version of  $\overline{M}(x)$ . If M(1) is non-zero, then since  $\overline{M}(1) = M(1)$ , we have that the polynomials are equal. If the degree of M(x) is even,  $\frac{M(x)}{x^{\deg(M)/2}}$  is a polynomial in  $x + \frac{1}{x}$ , contradicting the hypothesis that Q(x) is irreducible. When deg(M) is odd, we observe that M(-1) = 0, meaning that  $Q_R(-1) = 0$ , and thus that  $Q(-1 - \frac{1}{1}) = 0$ , meaning Q is of the form k(x + 2), for which the hypothesis is clearly true.

If instead M(1) = 0, since M(x) divides  $Q_R(x)$  and  $Q_R(1) = 0$ , we have  $Q(1 + \frac{1}{1}) = 0$ , and thus Q(x) = c(x - 2), which satisfies the condition.

2. If  $\overline{M}(x)$  divides N(x), and if  $\frac{N(x)}{\overline{M}(x)}$  has degree at least 1, we observe that  $M(x)M\left(\frac{1}{x}\right)$  is symmetric in  $\frac{1}{x}$ , which again leads to a contradiction with the assumption that Q(x) is irreducible. If no contradiction arises, then  $Q_R(x)$  must have exactly two irreducible factors, as needed.

Finally, combining these two facts, we consider an irreducible factor F(x) of  $P_n(x)$ , where n is as described in the first subclaim. We conclude that

$$x^{\deg(F)}F\left(x+\frac{1}{x}\right)$$

has degree  $2 \deg(F)$  and at most two prime factors of degrees at most Cn, completing the proof of the claim.

Now, consider n such that all irreducible factors of  $P_n(x)$  have degree at most  $\frac{n}{2024^{2025}}$ . We will show that such n satisfy the condition. Indeed, take an irreducible factor F(x) of the polynomial, and assume  $p \mid F(k)$  for a natural number k. By Bézout's Lemma, for any two coprime polynomial factors  $F_1(x)$  and  $F_2(x)$  of  $P_n$  (since they are irreducible and distinct), there exist polynomials A(x) and B(x) with integer coefficients such that

$$F_1(x)A(x) + F_2(x)B(x) = c$$

where c is a constant. Here, it is important to notice that  $P_n(x)$  does not have repeated factors, as that would imply  $x^{6n} - 1$  also having double factors. Doing this for every pair of irreducible divisors, we obtain the existence of a constant M (independent of p) such that for any k,

$$v_p(P_n) \le v_p(F(k)) + M.$$

For an irreducible factor F(x), this is at most  $\log_p(F(k)) + M$ . Using the fact that for large enough  $x, F(x) < x^{\frac{n}{2 \cdot 2024} 2024}$ , and that  $\log_p(x) \leq \log_2(x)$ , we arrive at the desired conclusion.