

# The Golden Digits National Contest

1ST EDITION, FEBRUARY 2024



**Problem 1.** Determine all functions  $f : \mathbb{R}_+ \rightarrow \mathbb{R}_+$  which satisfy

$$f\left(\frac{y}{f(x)}\right) + x = f(xy) + f(f(x)),$$

for any positive real numbers  $x$  and  $y$ .

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**Solution.** Rewrite the condition in the statement using the map  $y \mapsto y/x$ . It then becomes

$$f\left(\frac{y}{xf(x)}\right) - f(y) = f(f(x)) - x. \quad (\dagger)$$

Let  $\lambda_x = 1/(xf(x))$ . It follows inductively that for any positive integer  $N$  we have

$$-\frac{f(y)}{N} < \frac{f(y\lambda_x^N) - f(y)}{N} = f(f(x)) - x = \frac{f(y) - f(y\lambda_x^{-N})}{N} < \frac{f(y)}{N}.$$

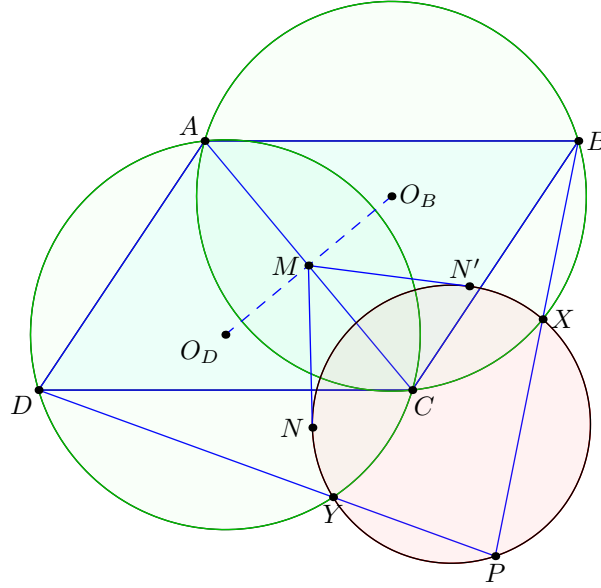
Taking  $N$  to be large enough, it follows that  $f(f(x)) = x$  so, in particular,  $f$  is injective. Recall the equation  $(\dagger)$ . These properties yield  $xf(x) = 1$  for every positive real number  $x$ .

To conclude, the only function which satisfies the condition is  $f(x) = 1/x$ , which works trivially.

**Problem 2.** Let  $ABCD$  be a parallelogram and  $P$  a point in the plane. The line  $BP$  intersects the circumcircle of  $ABC$  again at  $X$  and the line  $DP$  intersects the circumcircle of  $DAC$  again at  $Y$ . Let  $M$  be the midpoint of  $AC$ . The point  $N$  lies on the circumcircle of  $PXY$  so that  $MN$  is a tangent to this circle. Prove that  $MN$  and  $AM$  have the same length.

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**Solution 1.** Let  $\omega$  be the circumcircle of  $PXY$  and  $\omega_B, \omega_D$  the circumcircles of  $ABC$  and  $ADC$  respectively, with centers  $O_B, O_D$  and equal radius  $R$ . Also, consider the linear function  $f(Q) = AQ^2 - \text{Pow}_\omega(Q)$ . We wish to show that  $f(M) = 0$ , i.e.  $f(B) + f(D) = 0$ .



Observe that

$$\begin{aligned} f(B) &= AB^2 - BX \cdot BP \\ &= AB^2 - (BP - PX) \cdot BP \\ &= AB^2 - BP^2 + PX \cdot PB \\ &= AB^2 - BP^2 + PO_B^2 - R^2, \end{aligned}$$

where the last equality is obtained from the power of  $P$  with respect to  $\omega_B$ . Therefore, we have  $f(B) + f(D) = AB^2 + AD^2 - BP^2 - DP^2 + PO_B^2 + PO_D^2 - 2R^2$ . Using the median formula

$$4PM^2 = 2(BP^2 + DP^2) - BD^2 = 2(PO_B^2 + PO_D^2) - O_B O_D^2,$$

which yields

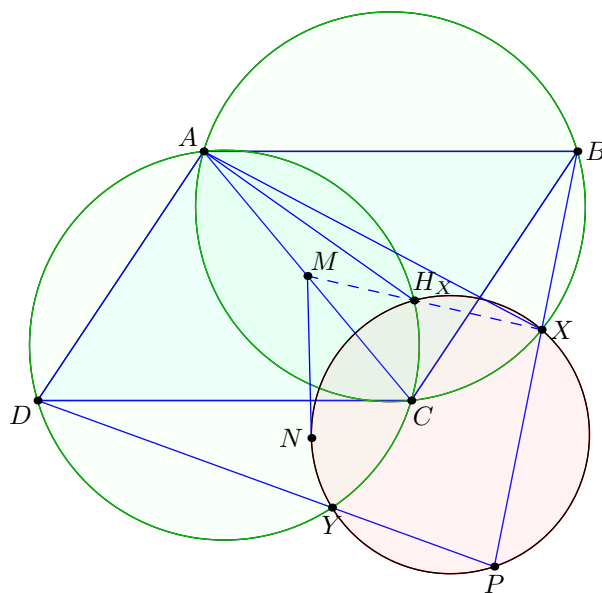
$$f(B) + f(D) = AB^2 + AD^2 + \frac{O_B O_D^2 - BD^2}{2} - 2R^2.$$

Finally, we have  $O_B O_D^2 = 4O_B M^2 = 4R^2 - AC^2$ , and  $2(AB^2 + AD^2) = AC^2 + BD^2$ , by the parallelogram formula. Combining these two, we get

$$f(B) + f(D) = AB^2 + AD^2 + \frac{4R^2 - AC^2 - BD^2}{2} - 2R^2 = 0,$$

exactly what we wanted to prove.

**Solution 2.** Consider  $H_X$ , the  $X$ -Humpty point in the triangle  $XAC$ . It is well-known that  $H_X$  lies on the median  $XM$ , as well as on the circle  $\omega_D$  (since  $\omega_B$  and  $\omega_D$  are symmetric with respect to  $AC$ ).



A quick angle chase gives

$$\begin{aligned} \angle H_X Y D &= \angle H_X A D = \angle H_X A C + \angle C A D \\ &= \angle H_X X A + \angle A C B \\ &= \angle H_X X A + \angle A X B \\ &= \angle H_X X B, \end{aligned}$$

or equivalently,  $\angle H_X Y P = \angle H_X X P$ , implying that  $H_X \in \omega$ . But then,

$$M N^2 = \text{Pow}_\omega(M) = M H_X \cdot M X = M A^2,$$

where the last equality is yet another well-known property. This concludes the solution.

**Problem 3.** There are  $m$  identical rectangular chocolate bars and  $n$  people. Each chocolate bar may be cut into two (possibly unequal) pieces at most once. For which  $m$  and  $n$  is it possible to split the chocolate evenly among all the people?

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**Solution.** It is possible if and only if  $m \geq n$  or  $m/n = 1 - 1/k$  for some integer  $k \geq 2$ .

We begin by showing that these values do work. Let  $\ell$  be the length of the chocolate bars. If  $m \geq n$  then place the chocolate bars in a line, forming a rectangle of length  $m\ell$ . Now, simply perform vertical cuts which split this large rectangle into  $n$  equal sections.



Because  $m \geq n$  the distance between any two cuts is  $m\ell/n \geq \ell$  hence each chocolate bar is cut at most once. Now, assume that  $m/n = 1 - 1/k$  for some integer  $k \geq 2$ . Observe that in this case,  $k - 1$  divides  $m$ . Split the chocolate bars into groups of size  $k - 1$ .

Like before, combine the chocolate bars in each group to form a rectangle of length  $\ell(k - 1)$  and perform vertical cuts which split each such rectangle into  $k$  equal sections. Thus, we divide the chocolate evenly into  $km/(k - 1) = n$  portions.

Note that we cut each chocolate bar at most once. Every chocolate piece created in this manner has length  $t \cdot \ell/k$  for some integer  $t \geq 1$ . Thus, if a chocolate bar would be cut at least twice, two of these cuts would be at a distance of at most  $(k - 2)\ell/k < (k - 1)\ell/k$ , absurd.

Assume that the condition holds for some  $m$  and  $n$  with  $m < n$ . We may also assume, for the sake of simplicity, that each chocolate bar has area one. Then, each person should receive a total area of chocolate equal to  $m/n < 1$  so each chocolate bar must be cut.

Suppose that after the cuts, the resulting chocolate pieces have areas  $a_1 \leq a_2 \leq \dots \leq a_{2m}$ . Of course, for each  $i$ , the pieces of areas  $a_i$  and  $a_{2m+1-i}$  make up a chocolate bar. We claim that  $a_i/a_1$  is always an integer. We proceed inductively; of course, this holds for  $i = 1$ .

Before continuing, observe that  $a_{2m} = m/n$ . If  $a_{2m} > m/n$  then this piece is too large by itself. If  $a_{2m} < m/n$ , then any other piece of chocolate contributes with an area of at least  $a_1$ , which is too much, as  $a_{2m} + a_1 = 1 > m/n$ . Either way, the chocolate cannot be split evenly.

Now, assume that for  $i = 1, \dots, k$  it is true that  $a_i/a_1$  is an integer. Note that since  $a_{k+1} + a_{2m-k} = 1 > m/n$  then  $a_{2m-k}$  adds up to  $m/n$  with some  $a_{i_1}, \dots, a_{i_l} < a_{k+1}$ . Hence, the indices  $i_1, \dots, i_l$  are at most  $k$  so due to the induction,  $(a_{i_1} + \dots + a_{i_l})/a_1$  is an integer.

Hence, we may write  $a_{2m-k} + t \cdot a_1 = m/n$ . Observe further that  $m/n = a_{2m} = 1 - a_1$  hence we may conclude that  $(t + 1) \cdot a_1 = 1 - a_{2m-k} = a_{k+1}$ . Therefore,  $a_{k+1}/a_1$  is an integer too, as desired. This completes the induction.

Now, we may write  $m/n = a_{2m} = (t - 1) \cdot a_1$  and  $1 = a_{2m} + a_1 = t \cdot a_1$  for some integer  $t \geq 2$  to infer that  $m/n = 1 - 1/t$ , as we wanted to prove. To conclude, the only  $m$  and  $n$  that work are the ones described above.